

CS 112: Computer System Modeling Fundamentals

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Lecture 18

Reminders & Announcements

- Homework 5 is due in section on Friday, which will be a review session for the final exam
- Check the website and catch up on your reading now!
- The best way to prepare for the final is to practice working through the problems in the book

Markov Chains

A Markov chain is specified by:

- A set of **states** $S = \{1, 2, \dots, m\}$
- A set of **transition probabilities** $p_{i,j}$ where

$$p_{i,j} = \mathbb{P}(X_{t+1} = j \mid X_t = i)$$

The key independence assumption is the **Markov property**:

$$\begin{aligned} \mathbb{P}(X_{t+1} = j \mid X_t = i, X_{t-1} = x_{t-1}, X_{t-2} = x_{t-2}, \dots, X_0 = x_0) \\ = \mathbb{P}(X_{t+1} = j \mid X_t = i) \end{aligned}$$

Classification of States

Accessibility: State j is **accessible** from state i if for some n , the n -step transition probability from i to j is positive.

Recurrence: State i is **recurrent** if for every state j that is accessible from i , i is also accessible from j .

If i is not recurrent, then it is **transient**.

The set of all states accessible from a recurrent state form a **recurrent class**. Note that all of the states in a recurrent class are accessible from each other.

Periodicity

- A recurrent class is **periodic** if it can be broken into $d > 1$ disjoint subsets S_1, S_2, \dots, S_d in such a way that
 - All transitions from states in S_i lead to states in S_{i+1} for $i \in \{1, \dots, d-1\}$
 - All transitions from states in S_d lead to states in S_1
- The **period** of the class is the number d of subsets

n-Step Transitions

- We can efficiently compute the *n*-step transition probability, $P(X_n = j \mid X_0 = i)$, using the recursive formula

$$P(X_n = j \mid X_0 = i) = \sum_{k=1}^m p_{k,j} P(X_{n-1} = k \mid X_0 = i)$$

Today...

- Long-term behavior of Markov chains
 - Steady state probabilities & the balance equations

(Break for course evaluations)

- Application: Google's PageRank algorithm

Back to the Faulty Router

My faulty router can be either online or offline. If it is online one day, it will be online the next day with probability 0.8. If it is offline one day, it will remain offline the next day with probability 0.4.

- What fraction of the time will my router be online in the long run?

Convergence of Markov Chains

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- Under what circumstances might this *not* happen?
 - Multiple recurrent classes
 - Periodic recurrent class
- If these probabilities *do* converge, what do we know about the values they converge to?

Steady-State Convergence Theorem

Theorem: Consider any Markov chain with a **single recurrent class**, which is **not periodic**. Then the **steady state probabilities** of the chain are the unique values π_1, \dots, π_m that satisfy the system of equations:

$$\pi_j = \sum_{k=1}^m \pi_k p_{k,j} \quad \text{for } j = 1, \dots, m$$
$$\sum_{k=1}^m \pi_k = 1$$

These are referred to as the **balance equations**.

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When these conditions hold, X_0 doesn't matter.

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- Suppose that if the router remains offline for 3 straight days, I (temporarily) repair it, resetting it to the online state. What fraction of the time will it be online?

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 - High quality webpages link to other high quality webpages. The rank of a webpage should be a function of the rank of pages that link to it.
 - If a webpage links to n other pages, each should inherit a $1/n$ share of its rank.

Google's PageRank

- Let S_i be the set of pages that link to page i , and let $n_j > 0$ be the number of pages that j links to. Then we want

$$R_i = \sum_{j \in S_i} R_j \frac{1}{n_j} \quad \text{for all pages } i$$

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- These equations can be interpreted as the balance equations of a **random surfer** Markov chain!
- Unfortunately, there might not be a unique solution...

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- This new MC has one recurrent class, and it is not periodic.

Other Applications

Similar ideas have been used in a variety of applications:

- Measuring the impact of bloggers in the blogosphere
- Measuring the impact of scientific journals based on citations
- Measuring how trustworthy buyers and sellers are on eBay or other e-commerce sites
- Determining the importance of species in the food chain